Chapter 6
Random Number Generation

- ❑ Requirements / application
- ❑ Pseudo-random bit generator
- ❑ Hardware and software solutions

Requirements and Application Scenarios

- ❑ Security
  - ❑ ***Key generation*** – automatically generated keys must be secure against "prediction" or "estimation"
  - ❑ ***Initialization vectors*** – many encryption algorithms rely on an IV, thus must be random to prevent guessing
  - ❑ ***Authentication*** – security protocols relying on challenge-response exchanges require random numbers
  - ❑ Further applications in cryptographic algorithms
- ❑ Other domains
  - ❑ Probabilistic decisions – if not "random", sequences may be created in long-term applications leading to self-similar behavior
  - ❑ Simulation techniques – calculation of variables following a particular distribution

Random Number Generation

- Hardware-based random bit generators are based on physical phenomena, as:
  - elapsed time between emission of particles during radioactive decay,
  - thermal noise from a semiconductor diode or resistor,
  - frequency instability of a free running oscillator,
  - the amount a metal insulator semiconductor capacitor is charged during a fixed period of time,
  - air turbulence within a sealed disk drive which causes random fluctuations in disk drive sector read latencies, and
  - sound from a microphone or video input from a camera
- A hardware-based random bit generator should ideally be enclosed in some tamper-resistant device and thus shielded from possible attackers

Random Number Generation

- Software-based random bit generators, are based upon processes as:
  - the system clock,
  - elapsed time between keystrokes or mouse movement,
  - content of input- / output buffers
  - user input, and
  - operating system values such as system load and network statistics
- Ideally, multiple sources of randomness should be “mixed”, e.g. by concatenating their values and computing a cryptographic hash value for the combined value, in order to avoid that an attacker might guess the random value
  - If, for example, only the system clock is used as a random source, than an attacker might guess random-numbers obtained from that source of randomness if he knows about when they were generated

Random Number Generation

- *De-skewing:*
  - Consider a random generator that produces biased but uncorrelated bits, e.g. it produces 1's with probability $p \neq 0.5$ and 0's with probability $1 - p$, where p is unknown but fixed
  - The following technique can be used to obtain a random sequence that is uncorrelated and unbiased:
    - The output sequence of the generator is grouped into pairs of bits
    - All pairs 00 and 11 are discarded
    - For each pair 10 the unbiased generator produces a 1 and for each pair 01 it produces a 0
  - Another practical (although not provable) de-skewing technique is to pass sequences whose bits are correlated or biased through a cryptographic hash function such as MD-5 or SHA-1

Statistical Tests for Random Numbers

- The following tests allow to check, if a generated random or pseudo-random sequence inhibits certain statistical properties:
  - *Monobit Test:* Are there equally many 1's like 0's?
  - *Serial Test (Two-Bit Test):* Are there equally many 00-, 01-, 10-, 11-pairs?
  - *Runs Test:* Are the numbers of *runs* (sequences containing only either 0's or 1's) of various lengths as expected for random numbers?
  - *Autocorrelation Test:* Are there correlations between the sequence and (non-cyclic) shifted versions of it?
  - *Maurer's Universal Test:* Can the sequence be compressed?
- The above descriptions just give the basic ideas of the tests. For a more detailed and mathematical treatment, please refer to sections 5.4.4 and 5.4.5 in [Men97a]

Random and Pseudo-Random Number Generation

- Definition: A ***random bit generator*** is a device or algorithm, which outputs a sequence of statistically independent and unbiased binary digits.
  - Remark: A random bit generator can be used to generate uniformly distributed random numbers, e.g. a random integer in the interval [0, n] can be obtained by generating a random bit sequence of length $\lfloor \lg n \rfloor + 1$ and converting it into a number. If the resulting integer exceeds *n* it can be discarded and the process is repeated until an integer in the desired range has been generated.

Random and Pseudo-Random Number Generation

- Definition: A ***pseudo-random bit generator (PRBG)*** is a deterministic algorithm which, given a truly random binary sequence of length *k*, outputs a binary sequence of length *m* >> *k* which "appears" to be random. The input to the PRBG is called the *seed* and the output is called a *pseudo-random bit sequence*.
  - Remarks:
    - The output of a PRBG is not random, in fact the number of possible output sequences of length *m* is at most all small fraction $2^k / 2^m$, as the PRBG produces always the same output sequence for one (fixed) seed
    - The motivation for using a PRBG is that it might be too expensive to produce true random numbers of length *m*, e.g. by coin flipping, so just a smaller amount of random bits is produced and then a pseudo-random bit sequence is produced out of the *k* truly random bits
    - In order to gain confidence in the "randomness" of a pseudo-random sequence, statistical tests are conducted on the produced sequences

Random and Pseudo-Random Number Generation

- Example:
  - A linear congruential generator produces a pseudo-random sequence of numbers $y_1, y_2, \ldots$ According to the linear recurrence

    $$y_i = a \times y_{i-1} + b \bmod q$$

    with a, b, q being parameters characterizing the PRBG
  - Unfortunately, this generator is predictable even when a, b and q are unknown, and should, therefore, not be used for cryptographic purposes

Random and Pseudo-Random Number Generation

- Security requirements of PRBGs for use in cryptography
  - As a minimum security requirement the length k of the seed to a PRBG should be large enough to make brute-force search over all seeds infeasible for an attacker
  - The output of a PRBG should be statistically indistinguishable from truly random sequences
  - The output bits should be unpredictable for an attacker with limited resources, if he does not know the seed

- Definition: A PRBG is said *to pass all polynomial-time statistical tests*, if no polynomial-time algorithm can correctly distinguish between an output sequence of the generator and a truly random sequence of the same length with probability significantly greater than 0.5
  - *Polynomial-time algorithm* means, that the running time of the algorithm is bound by a polynomial in the length m of the sequence
- Definition: A PRBG is said *to pass the next-bit test*, if there is no polynomial-time algorithm which, on input of the first m bits of an output sequence s, can predict the $(m + 1)^{st}$ bit s_{m+1} of the output sequence with probability significantly greater than 0.5
- Theorem (universality of the next-bit test):

  A PRBG passes the next-bit test $\Leftrightarrow$ it passes all polynomial-time statistical tests
  - For the proof, please see section 12.2 in [Sti95a]

- Definition: A PRBG that passes the next-bit test – possibly under some plausible but unproved mathematical assumption such as the intractability of the factoring problem for large integers – is called a ***cryptographically secure pseudo-random bit generator (CSPRBG)***

Pseudo-Random Number Generation

- There are a number of algorithms, that use cryptographic hash functions or encryption algorithms for generation of cryptographically secure pseudo random numbers
  - Although these schemes can not be proven to be secure, they seem sufficient for most practical situations
- One such approach is the ANSI X9.17 generator:
  - Input: a random and secret 64-bit seed s, integer m, and 3-DES key K
  - Output: m pseudo-random 64-bit strings $y_1, y_2, \ldots Y_m$

    1.) $q = E(K, Date\_Time)$

    2.) For i from 1 to m do

    2.1) $x_i = E(K, (q \oplus s)$

    2.2) $s = E(K, (x_i \oplus q)$

    3.) Return($x_1, x_2, \ldots x_m$)
  - This method is a U.S. Federal Information Processing Standard (FIPS) approved method for pseudo-randomly generating keys and initialization vectors for use with DES

Secure Pseudo-Random Number Generation

- The RSA-PRBG is a CSPRBG under the assumption that the RSA problem is intractable:
  - Output: a pseudo-random bit sequence $z_1, z_2, \ldots, z_k$ of length k

    1.) Setup procedure:

    Generate two secret primes p, q suitable for use with RSA
    Compute $n = p \times q$ and $\Phi = (p - 1) \times (q - 1)$
    Select a random integer e such that $1 < e < \Phi$ and $\gcd(e, \Phi) = 1$

    2.) Select a random integer y_0 (the seed) such that $y_0 \in [1, n]$

    3.) For i from 1 to k do

    3.1) $y_i = (y_{i-1})^e \bmod n$

    3.2) z_i = the least significant bit of y_i
  - The efficiency of the generator can be slightly improved by taking the last j bits of every y_i, with $j = c \times \lg(\lg(n))$ and c is a constant
  - However, for a given bit-length m of n, a range of values for the constant c such that the algorithm still yields a CSPRBG has not yet been determined

Secure Pseudo-Random Number Generation

- The Blum-Blum-Shub-PRBG (BBS) is a CSPRBG under the assumption that the integer factorization problem is intractable:
  - Output: a pseudo-random bit sequence $z_1, z_2, ..., z_k$ of length k

    1.) Setup procedure:

    Generate two large secret and distinct primes p, q such that p, q are each congruent 3 modulo 4 and let $n = p \times q$

    2.) Select a random integer s (the seed) such that $s \in [1, n - 1]$ such that gcd(s, n) = 1 and let $y_0 = s^2 \bmod n$

    3.) For i from 1 to k do

    3.1) $y_i = (y_{i-1})^2 \bmod n$

    3.2) z_i = the least significant bit of y_i
  - The efficiency of the generator can be improved using the same method as for the RSA generator with similar constraints on the constant c

Summary (what do I need to know)

- Principles
  - Random bit generator
  - Pseudo-random bit generator
  - Cryptographically secure pseudo-random bit generator
- Hardware solutions
  - Examples
- Software solutions
  - Examples